## Ch 5 - More Prolog features

## Towards « real » programming

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## Remember: arithmetics

In order to ask Prolog to treat numbers, use is/2 (that is extra-logical)
? $-X=2+3$.
$X=2+3$
?- X is $2+3$
$X=5$
?- X is $\mathrm{Y}+3$
Error

## Modes

- Not all predicates are fully declarative
- It is important to know the mode of the arguments when the goal is called
- ++: should be ground
-     + : should not be a variable (but can contain variable)
-     - : should be a variable
- ? : can be not instantiated at all

Crucial information in the library documentation

- Example
- factorial(++, ?)
- cannot be used with variables in the first argument


## Exercise 5.1: Zebra puzzle $1 / 2$

The "Zebra puzzle":
1 There are 5 colored houses in a row, each having an owner, which has an animal, a favorite cigarette, a favorite drink.
2 The English lives in the red house.
3 The Spanish has a dog.
4 They drink coffee in the green house
5 The Ukrainian drinks tea.
6 The green house is next to the white house.
7 The Winston smoker has a serpent.
8 In the yellow house they smoke Kool.
9 In the middle house they drink milk.
10 The Norwegian lives in the first house from the left.
11 The Chesterfield smoker lives near the man with the fox.
12 In the house near the house with the horse they smoke Kool.
13 The Lucky Strike smoker drinks juice.
14 The Japanese smokes Kent.
15 The Norwegian lives near the blue house.

Who has a zebra and who drinks water?

## Exercise 5.1: Zebra puzzle $2 / 2$

- Write a Prolog program to solve The Zebra problem
- The main predicate has 17 subgoals.
- How to proceed
- Represent the houses as a list with 5 lists from left to right in the street:

```
    Sol = [[Man1, Animal1, Cigarette1, Drink1, Color1],
        [..],..],[..],
        [Man5, Animal5, Cigarette5, Drink5, Color5] ]
```

- Define predicate $\operatorname{right}(\mathbf{X}, \mathbf{Y}, \mathbf{L})$ that is true if $X$ is just after $Y$ in list L .
- Define predicate near( $\mathbf{X}, \mathbf{Y}, \mathbf{L}$ ) that is true if either X is just after Y or Y is just after X is L .
- Use predicates member/2
- Test case : ?- zebra(Sol).



# Take your time to search, code and test your own program 

Then take your time to understand the following solution

## Exercise 5.1: Zebra puzzle 2/2 (bis)

zebra(Sol):-
> length(Sol, 5),
> member([english,_,_,_red], Sol),
> member([spanish,dog,_,_,], Sol),
> member([_,_,_,coffee,green], Sol),
> member([ukrainian,_,_,tea,_], Sol),
> \% 1
> right([_,_,_,_,green],[_,_,_,_, white], Sol), \% 6
> member([_,snake,winston,_,_], Sol), \% 7
> member([_,_,kool,_,yellow], Sol), \% 8
> Sol= [_,_,[_,_,_,milk,_],_,_], \% 9
> Sol= [[norwegian,_,_,_,_],_,_,_,_], \% 10

near([_,_,chesterfield,_,_],[_fox,_,_,_], Sol), \% 11
near([_,_,kool,_,_],[_,horse,_,_,_], Sol), \% 12
member([_,_lucky,juice,_], Sol), \% 13
member([japonese,_,kent,_,_], Sol), \% 14
near([norwegian,_,_,_,_],[_,_,_,_, blue], Sol), \% 15
member([_,_,_,water,_], Sol), \% someone drinks water
member([_,zebra,_,_,_], Sol). \% someone has a zebra
$\operatorname{right}\left(X, Y,\left[Y, X \mid \_\right]\right)$.
$\operatorname{right}\left(X, Y,\left[\_\mid Z s\right]\right):-$
$\operatorname{right}(X, Y, Z s)$.
left(X, Y, L) :-
$\operatorname{right}(Y, X, L)$.

## MORE IMPORTANT FEATURES

## More important features

- Input/output
- read/2, read/3
- printf/2, printf/3
- Controlling backtracking
- !/0 (cut)
- Negation as failure
- not/1
- Operators

Note that there are many more built-in predicates. See documentation: http://eclipseclp.org/doc/bips/

## Input: read(-Term)

Succeeds if the next term from the input stream is successfully read and unified with Term.
?- read(X).
12654.

X $=12654$
?- read( X ).
hello.
X = hello
?- read $(X)$.
father_of(Iali, ana).
X = father_of(Iali, ana).
?- read $(X)$.
father_of(lali.

## Error

read/1 is a full parser !!

It reads what is typed in the input stream

It builds a Prolog term of any complexity
except if there are syntax errors

## Input: read(-Term, ++Stream)

- read/2 behaves like read/1 but it reads from a given stream
- very useful if you want to read from a file
- In that case the programming pattern is top(FileName) :-
open(FileName, read, s),
my_program(..., s),
close(s).
my_program(...., s) :-
read(s, A),
do_something(A, ...),

See http://eclipseclp.org/doc/bips/kernel/ioterm/read-2.html

## Output: printf(+Format, ?ArgList)

- The arguments in the argument list ArgList are interpreted according to the Format string and the result is printed to the output stream
- A useful example
?- printf("\tHello \%w !\n\tYes, \%w !!", [you, 'I mean you']). Hello you!
Yes, I mean you !!


## output: printf(+Stream, +Format, ?ArgList)

- Same as printf/2 but can write on any file
- In that case the programming pattern is
top(FileName) :-
open(FileName, write, s),
my_program(..., s),
close(s).
my_program(...., s) :-
printf( "...", [...]),
- See the Eclipse documentation for the details of the format
http://eclipseclp.org/doc/bips/kernel/ioterm/printf-3.html


## Back to ancestor

- Add a predicate ancestor/0, that
- asks the user for whom s.he wants to find ancestors
- prints, for each solution, the ancestor

Do not forget that when you enter the person name you must end up with a '.'


# Take your time to search, code and test your own program 

Then take your time to understand the following solution

## Back to ancestor (bis)

- Add a predicate ancestor/0, that
- asks the user for whom s.he wants to find ancestors and
- prints, for each solution, the ancestor
- Code
ancestor :-
printf("Initial child?", []), read(C),
ancestor(A, C),
printf("\%w is an ancestor of: \%w w ", $[\mathrm{A}, \mathrm{C}]$ ).

Note that this is what you are used to with procedural programming languages but it is much less general than ancestor/2...

## Why?

## Controlling backtracking: !/0

- Backtracking is very powerful but sometimes we need to control it
- built-in predicate '!' (called 'cut') is used to tell the interpreter not to backtrack
- it is always true and works by side effects on the interpreter internal (hidden) structure
- it cuts branches in the search tree
- within a certain scope
- this is very useful but extra-logical


## cut : example 1

$$
\begin{gathered}
\text { aa }:- \\
\text { bb1, } \\
\text { bb2, } \\
!,
\end{gathered}
$$



If! is executed
other possibilities are abandoned
bb3.

all possibilities are tried

```
aa :-
    cc1,
        cc2.
```

bb1 :- dd1.
bb1 :- dd2.
bb2 :- ee1.
bb3 :- ff1.
bb3 :- ff2.

Tried
dd1. dd2.
ff1.
ff2.

## Cut: example 2 1/3

pp0(X, Y) :-

```
        qq(X),
    qq(Y).
pp0(0,1).
qq(1). qq(2).
```



$$
\mathrm{X}=0, \mathrm{Y}=1
$$

## Cut: example2 2/3

ppl(X, Y) :-
$q q(X)$. $q q(Y)$,
!.
$\operatorname{ppl}(0,1)$.
qq(1).
$q q(2)$.

$\mathrm{qq}(\mathrm{Y})$, !

succès
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## Cut: example2 3/3



## Back to

french_menu/2

```
french_menu([A, M], Cal) :-
        appetizer(A,ApCal),
    main_course(M , MaCal),
    check_cal([ApCal, MaCal], Cal).
    french_menu([M, D], Cal) :-
    main_course(M , MaCal),
    dessert_or_cheese(D, DeCal),
    check_cal([MaCal, DeCal], Cal).
```

What are the answers to the following queries ? ?- french_menu(M, C), !.
?- french_menu([X, Y], C), main_course(X, _), !.
?- french_menu([X, Y], C) , !, main_course(X, _).

## Back to ex 3.5 deleteXs(X, L1, L2)

We designed deleteXs(_X, [], []). deleteXs(X, [X|L1], L2) :deleteXs(X, L1, L2). deleteXs(X, [Z | L1], [Z | L2]) :-
$X 1=Z$,
deleteXs(X, L1, L2).
How could we prune the search tree without loosing any solution?


# Take your time to search, code and test your own program 

Then take your time to understand the following solution

## Back to ex 3.5 deleteXs/3: update 1 (bis)

We designed
deleteXs(_X, [], []).
deleteXs(X, [X|L1], L2) :deleteXs(X, L1, L2).
deleteXs(X, [Z | L1], [Z | L2]) :-
$X l=Z$, deleteXs(X, L1, L2).
How could we prune the search tree Note that the base clause without loosing any solution?
deleteXs(_X, [], []).
deleteXs(X, [X | L1], L2) :-
deleteXs(X, L1, L2),
!.
deleteXs(X, [Z | L1], [Z | L2]) :\% $\quad \mathrm{X}=\mathrm{Z}$, deleteXs(X, L1, L2).
$\% . .$. <end of line>:
comment does not need a cut because Prolog compiler is clever enough to deduce that it is exclusive

The test is not necessary, but keeping it would be correct

## Cut and declarative programming

- Cut is extra-logical
- testing all cases is especially crucial
- remember that you should always test
- verification of at least a correct solution (that should get 'yes')
- verification of at least an incorrect solution (that should get ' $n o$ ')
- generations of solutions (at least one test case per argument with that argument variable)


## Exercise 5.2: minimum of 2 integers

- $\min (M, X, Y)$ is true if $M$ is the minimum of $X$ and Y
- Write 2 versions
- one without cut and one with cut that prunes the search tree without changing the results
- specify the mode of the arguments when the goal is called
- ++: should be ground
-     + : should not be a variable (but can contain variable)
-     - : should be a variable
- ? : can be not instantiated at all



# Take your time to search, code and test your own program 

Then take your time to understand the following solution

## Ex. 5.2: minimum of 2 integers (bis)

- $\min (M, X, Y)$ is true if $M$ is the minimum of $X$ and Y
- Write 2 versions
- one without cut and one with cut that prunes the search tree without changing the results
:- mode mini(?,,++++ ).
$\operatorname{mini}(X, Y, X)$ :-

$$
X<Y .
$$

$\operatorname{mini}(X, Y, Y):-$
$Y=<X$.
Optimized version $\operatorname{mini}(X, Y, X)$ :-
$X<Y$, !.
$\operatorname{mini}(X, Y, Y):-$
$Y=<X$.

## Ex. 5.2: minimum of 2 integers (ter)

This is incorrect $\operatorname{mini}(X, Y, X):-$

$$
X<Y,
$$

$$
!.
$$

## ?- $\min (2,5,5)$. yes

(It should be "No", 2 is not the minimum of 5 and 5 !)

## Assumption of a closed world

- Negation as failure
- if it cannot be proved $=$ it is considered negated
- The standard procedure is to check if a goal succeeds
- You can explicitly check if a goal fails
- using predicate not/1
- But you have to be careful


## Predicate not/1

- Existing built-in meta-predicate
- namely a predicate that takes a predicate as argument
not(P) :-
P,
!, fail.
not(P).
fail/0 is another built-in predicate
It forces the execution to fail.

The second clause is only tried if $P$ did not previous succeed, hence telling that $P$ fails

## Exercise 5.3: small/1

short(X) :not(tall(X)).
tall(peter).
tall(paul).

- What are the answers to
?-short(mary).
?-short(peter).
?-short(X), X=mary.



# Take your time to search, code and test your own program 

Then take your time to understand the following solution

## Exercise 5.3: small/1 (bis)

short(X) :not(tall(X)).
tall(peter).
tall(paul).

When using not/1 you
should be careful about non ground arguments !
! not/1 is not fully logical

- What are the answers to
?-short(mary).
?-short(peter).
?-short(X), X=mary.
yes (valid)
no (valid)
no (invalid)


## Operators

Operators help to write more readable code

- useful when your code is to be read by non experts
you have to declare
- priority
- whether is an infix, prefix or suffix operator
- (sometimes tricky)

Examples

- $X$ parent_of $Y$ is the same as parent_of $(X, Y)$
:- op(500, xfx , parent_of).
infix operator with priority 500
- X is $3+2$ is the same as is $(\mathrm{X}, 3+2)$

Note that the project does not need operators

## ECLiPSe Documentation

- User manual
- Tutorials
- The Reference Manual(s)
- must always be open while programming
- predicate with the largest arity often the most general one, with the most detailed help
- eg min_max/8
- On line help
- same contents as reference manual
:- help <keyword>.
http://eclipseclp.org/doc/index.html


## ECLiPSe Documentation

- ECLiPSe Tutorial Introduction, also in pdf format
- Developing Applications with ECLiPSe, also in pdf format
- User Manual, also in pdf format
- Constraint Library Manual, also in pdf format
- Reference Manual (Built-In Predicates and Libraries) with Alphabetical Predicate Index
- Embedding and Interfacing Manual, also in pdf format
- API documentation for the Java-Eclipse Interface
- Visualisation tools manual, also in pdf format
- Obsolete Libraries Manual, also in pdf format
- Constraint Programming Examples (ECLiPSe web site)
- Examples for Embedding (C, C++, VBasic, Java) and Search
- ECLiPSe web site
- How to report a bug
- Join the mailing list!

Third party components:

- $\operatorname{Clp}(\mathrm{Q}, \mathrm{R})$ Library Manual (Postscript)


## ECLiPSe 6.0 Reference Manual

1. The ECLiPSe Built-In Predicates
2. The ECLiPSe Libraries
3. Third Party Libraries

## Built-Ins and Libraries by Categories

## Built-In Predicates

allsols arithmeticl compilerl controll debugl directives dynamicl envl eventl externalsl iocharl iostreaml ioterml modulesl obsoletel opsys l record storagel stratoml suspensionsl syntax l termcompl termmanipl typetest
Algorithms
all_min_cutsl all_min_cuts_eplexl anti_ unifyl applyl apply_macrosl bfsl branch and boundl calendarl changeset colgenl edge finderl edge finder3l fd global gacl graph algorithms ic global gacl max _low max flow_eplex notinstancel numbervars par_utill regex suspendl tentative_constraints|

## Compatibility

attsl ciol conjunto _fd setsl cprolog f fompilel foreign isol mercuryl multifile numbervars obsoletel quintus sepial sicstus socketss swil

## Constraints

bfsl cardinall changesetl chrl colgenl conjuntol conjunto fd setsl constraint poolsl cumulativel cyclel ech edge finderl edge finder3l eplex eplex cplex eplex osil eplex osi_clpcbcl eplex osi_ symclpl
 gras erl icl ic uumulativel ic edge finderl ic edge finder3l ic _gap_sbdd ic_gap_sbds ic globall
ic_s han and ic hybrid setss ic kernell ic make_overlap.
ic probe_support ic_ probing_for_scheduling ic sbds ic its ic symbolic 1 sbl make_overlap_ bivs minizincl mipl probel probe_searchl probe_supportl probing E. .ocodolved. opial repairl sdl shadow_conss suspend sym_ exprI tentativel tentative_constraints|
Data Structures
config_optss constraint_pools graph algorithms hashl heapsl linearizel list_collectil nl lists| 1 tutl 1 m mapl
m _tree2341 matrix utill notify_ports ordsetl queuess record shadow_consl storagel ar name
Development Tools
asml compilerl coveragel debugl documentl env| fcompilel instprofilel instrument lintl lips mode_ analyser| port _profilerl pretty_printl pretty_printer profile remote_tools source_ processorl spelll test_utill time_log toplevell vc_support viewablel xrefl

## flatten(+NestedList, ?FlatList)

Succeeds if FlatList is the list of all elements in NestedList, as found in a left-to-right, depth-first traversal of NestedList,

+ NestedList
Ground List.
?FlatList
List or variable.


## Description

FlatList is the list built from all the non-list elements of NestedList and the flattened sublists. The sequence of elements in FlatList is determined by a eft-to-right, depth-first traversal of NestedList.

The definition of this Prolog library predicate is:
flatten(List, Flat) :-
flatten_aux(List, Flat, []).
flatten_aux([], Res, Cont) :- -?-> !, Res = Cont.
flatten_aux([Head|Tail], Res, Cont) :-
-?->
flatten_aux(Head, Res, Cont1),
flatten_aux (Tail, Cont1, Cont)
This predicate does not perform any type testing functions.

## Modes and Determinism

- flatten $(+,-)$ is det


## Fail Conditions

Fails if FlatList does not unify with the flattened version of NestedList.

## Resatisfiable

No.

## Examples

Success:
[eclipse]: flatten([[1,2,[3,4],5],6,[7]], L).
$L=[1,2,3,4,5,6,7]$
yes.
Fail:
[eclipse]: flatten([1,[3],2], [1,2,3])
no.

## See Also

